

# Schedule-Free High-Frequency Transit Operations

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Gabriel E. Sánchez-Martínez

Nigel H.M. Wilson

Haris N. Koutsopoulos

Massachusetts Institute of Technology

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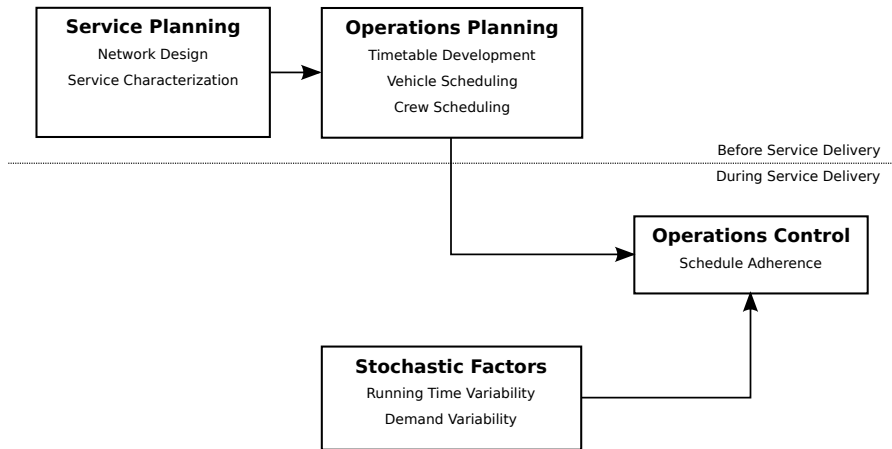
# Outline

- ① Framework
- ② Methodology
- ③ Application
- ④ Concluding Remarks

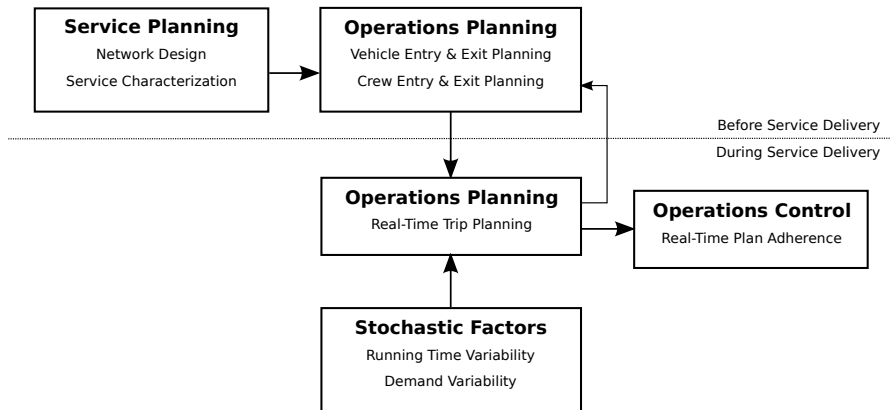
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# Schedule-Based Paradigm



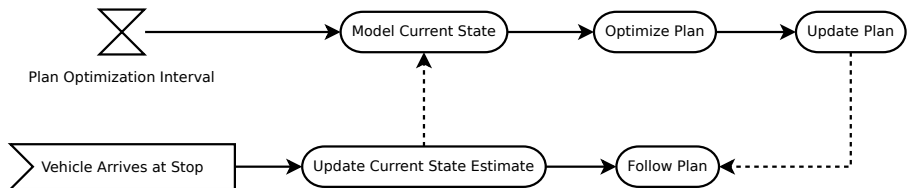
# Schedule-Free Paradigm



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# General Methodology



# Optimization Problem

$$\underset{x \in X}{\text{minimize}} \quad C(x; p)$$

- $C(x; p)$  is a general non-convex cost function covering the horizon
  - passenger cost: waiting and in-vehicle time
  - driver exit lateness cost
  - plan complexity cost
- $x$  is a candidate plan;  $X$  is the set of all feasible plans
- $p$  is a set of exogenous parameters and initial conditions
- Constraints
  - vehicles may not exit later than latest allowed exit time
  - vehicle must exit at required exit locations
  - lower and upper bounds on holding times



# Spatiotemporal Decomposition

Trip sequences are discrete, spatial variables

$$\underset{s \in S}{\text{minimize}} \quad C(s; d_s^*, p)$$

- $s$  is a candidate combination of trip sequences for all vehicles
- $S$  is the set of all feasible trip sequence combinations
- $d_s^*$  are optimal departure times for each given  $s$

Departure times are continuous, temporal variables

$$\underset{d \in D}{\text{minimize}} \quad C(d; s, p)$$

- $d$  is a set of departure times for all vehicles
- $D$  denotes the feasible space of departure times

# Problem Complexity

$$\mathcal{O} \left( \prod_{v \in V} |S_v| \sum_{i=1} k_i \right)$$

- $S_v$  is the set of candidate trip sequences for vehicle  $v$
- $k_i$  is the complexity of optimizing departure times for a set of trip sequences  $i$

Example: 20 vehicles, 5 sequences per vehicle, constant complexity  $k$

$$\mathcal{O}(k \cdot 5^{20})$$

# Simplified Methodology

Decomposition into sequentially solved subproblems for each vehicle

$$\underset{s_v \in S_v}{\text{minimize}} \quad C(s_v; s_{\bar{v}}, d_s^*, p)$$

- $s_v$  is a candidate trip sequence;  $S_v$  is the set of feasible trip sequences
- $s_{\bar{v}}$  are the trip sequences assumed for all other vehicles
- $d_s^*$  are the optimal departure times for each trip sequence combination

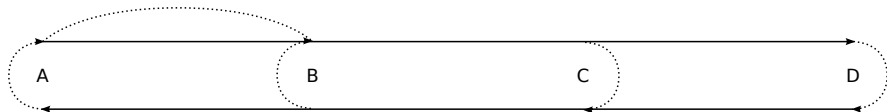
Example: 20 vehicles, 5 sequences per vehicle, constant complexity  $k$ ,  
 $g$  passes

$$\mathcal{O}(gk \cdot 5 \cdot 20)$$

# Optimization Algorithm

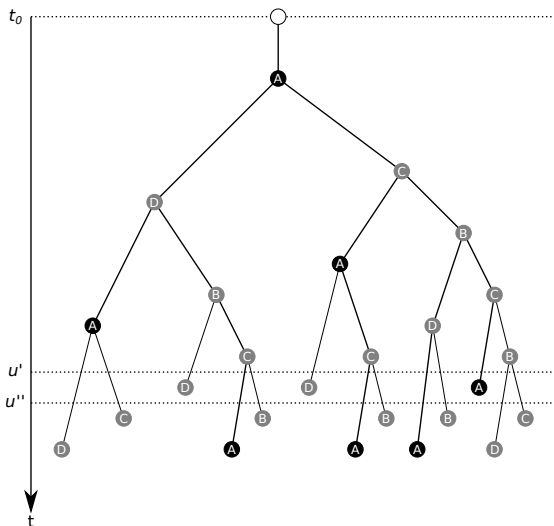
- ① Generate *basic trip sequences* for each vehicle
- ② Rank vehicles
- ③ Optimize trip sequence and departure times for each vehicle
  - ① Enumerate feasible trip sequences
  - ② Discard likely suboptimal trip sequences
  - ③ Optimize departure times for each trip sequence
    - constrained even headway algorithm
    - constrained rolling horizon optimization
  - ④ Evaluate performance for each trip sequence
  - ⑤ Select best plan
- ④ Update operations plan

# Modeling Variations



- trips between terminals
- short-turning (benign or aggressive)
- deadheading
- expressing (benign or aggressive, including skip-stop)

# Generating Feasible Trip Sequences



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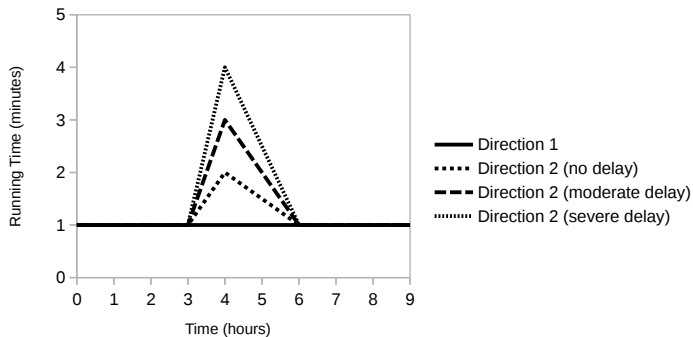
# Application



- 25 vehicles (not all operating simultaneously)
- each with capacity for 60 passengers
- trips dispatched over a period of 8 hours
- deterministic running times: no delay, moderate delay, severe delay
- Poisson demand, but with common random numbers
- vehicles should not exit more than 15 minutes after the last scheduled stop



# Application: Running Time Delays



- The schedule assumes no delays
- The real-time operations planning algorithm assumes no further delays

# Results: No Delay

Delay	Performance Measure	Short-Turning		No Short-Turning	
		SB	SF	SB	SF
None	Waiting Time (min)	2.6	2.6	2.6	2.6
	Excess Waiting Time (min)	0.0	0.0	0.0	0.0
	In-Vehicle Time (min)	9.6	9.5	9.5	9.5
	Late Exits	0	0	0	0
	Max Exit Lateness (min)	0.0	0.0	0.0	0.0
	Trips	190	192	190	190
	Short Turns	0	3	—	—

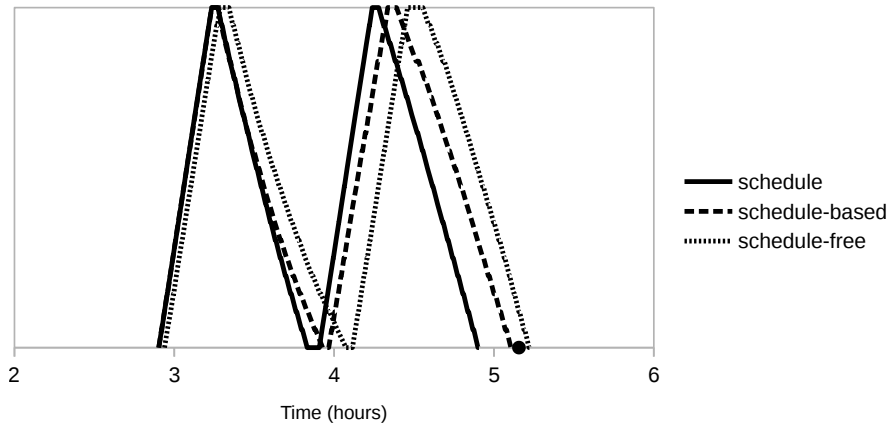
- schedule-based and schedule-free paradigms perform similarly

## Results: Moderate Delay

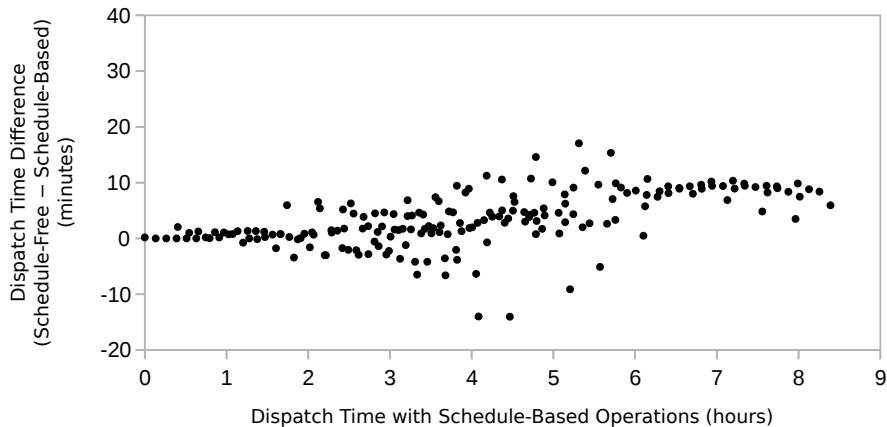
Delay	Performance Measure	Short-Turning		No Short-Turning	
		SB	SF	SB	SF
Moderate	Waiting Time (min)	3.3	2.7	3.4	2.7
	Excess Waiting Time (min)	0.7	0.1	0.8	0.2
	In-Vehicle Time (min)	10.8	10.7	10.7	10.7
	Late Exits	0	6	1	5
	Max Exit Lateness (min)	0.0	4.1	0.6	1.0
	Trips	190	190	190	186
	Short Turns	2	2	—	—

- schedule-free paradigm improves passenger experience
- schedule-free paradigm can lead to greater driver exit lateness
- allowance of short-turning can be counter-productive

# Trajectory of Vehicle with Latest Exit: Moderate Delay



# Schedule-free dispatch lateness



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# Concluding Remarks

- ① Under the SFP, vehicle trips are planned in real-time, utilizing real-time information.
- ② The real-time planning problem is formulated as a cost minimization problem, but decomposed into related subproblems for each vehicle for tractability.
- ③ The SFP is feasible and potentially beneficial. However, it can lead to greater exit lateness when delays are unexpected.

# Future Work

- 1 Stochasticity
- 2 Information modeling
- 3 Control strategies
- 4 Entry and exit plan optimization
- 5 Real-time planning optimization methods
- 6 Driver constraints
- 7 Autonomous fleets
- 8 Organizing informal systems
- 9 Policy implications
- 10 Tests on real transit lines



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